# **Connection Between Regular Expressions and Regular Languages** Section 3.2 [Textbook-1]

• for every regular language there is a regular expression

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• Regular Expressions denote Regular Languages

$$-\begin{array}{c} - \begin{array}{c} q_{0} \\ q_{1} \end{array} \end{array} \\ (a) \end{array} \\ (b) \end{array} \\ (b) \end{array} \\ (c) \end{array}$$

(a) nfa accepts Ø.

(b) nfa accepts  $\{\lambda\}$ .

(c) nfa accepts  $\{a\}$ .



#### Schematic representation of an nfa accepting L(r).



#### Automaton for $L(r_1 + r_2)$ .



#### Automaton for $L(r_1r_2)$ .





#### Exercise

Find an nfa that accepts L(r), where

 $r=(a+bb)*(ba*+\lambda)$ 

for every regular language, there should exist a corresponding regular expression.

- For every regular language there should be an NFA that accepts it.
- From the NFA, we can extract the respective regular expression using  $g_{\rm em}^{M}$  and  $g_{\rm em}^{M}$  an

From construct the equivalent Generalized Transition Graph in which transition labels are regular expressions.



# **Reducing the states**

**Reducing the states** 



# Example (cont.)

**Resulting Regular Expression:** 



# r = (bb\*a)\*bb\*(a+b)b\*

$$L(r) = L(M) = L$$

# **In General**

**Removing states:** 



# **In General**

#### The final transition graph:



The resulting regular expression:

$$r = r_1 * r_2 (r_4 + r_3 r_1 * r_2) *$$
$$L(r) = L(M) = L$$

# Example

Find a regular expression for the language

$$L = \{w \in \{a, b\}^* : n_a(w) \text{ is even and } n_b(w) \text{ is odd} \}.$$

Steps:

- Build NFA that accepts this language
- Generate the complete GTG by adding edges between each pair of states in the NFA
- Use the described state reduction procedure to reach the final transition graph, then apply the equation in the previous slide to find the regular expression.
- [Note: in Section 2.3, students can find a detailed procedure on this **nfa-to-rex** process]