

# Key Exchange Solutions

- Diffie-Hellman Protocol
- Needham Schroeder Protocol
- X.509 Certification

# Diffie-Hellman Key Exchange

- The Diffie-Hellman protocol allows 2 people to use random values and yet each generate the same symmetric key without transmitting the value of the key.
- The security of the protocol lies in the discrete log problem:

given  $y$ ,  $g$  and  $p$  find  $x$  such that

$$y = g^x \text{ mod } p$$

Alice and Bob need to agree on a key to use in a symmetric key cryptosystem. They choose a large prime number  $p$  and generator  $g$ .

### Alice

### Bob

1. Generates random number  $a$ ,
2. Computes  $x=g^a \bmod p$
3. Sends  $x$  to Bob
4. Receives  $y$  from Bob
5. Computes  $k=y^a \bmod p$

1. Generates random number  $b$ ,
2. Computes  $y=g^b \bmod p$
3. Sends  $y$  to Alice
4. Receives  $x$  from Alice
5. Computes  $k=x^b \bmod p$

# Why Diffie-Hellman works

Alice has computed

$$\begin{aligned}k &= y^a \pmod p \\ &= (g^b)^a \pmod p \\ &= g^{ba} \pmod p \\ &= g^{ab} \pmod p\end{aligned}$$

Bob has computed

$$\begin{aligned}k &= x^b \pmod p \\ &= (g^a)^b \pmod p \\ &= g^{ab} \pmod p\end{aligned}$$

So Alice and Bob both have the same value of  $k$ .

# How secure is it?

We assume that cryptanalyst Charles knows the values of  $p$  and  $g$  and that he eavesdrops on the exchange between Alice and Bob so that he also knows  $x$  and  $y$ .

However, unless Charles can solve a DLP, he is unable to find  $a$  or  $b$ .

It is believed that it is just as hard to find  $k$  from  $x$  and  $y$  without finding  $a$  or  $b$ .

# The Needham-Schroeder Protocol

This is another protocol for exchanging keys between Alice and Bob.

This time they use only symmetric key cryptography

But

They need a trusted third party (TTP) or Server (S).

- Alice and the server have a key  $K_{AS}$
- Bob and the server have a key  $K_{BS}$
- Alice and Bob want to establish a shared key  $K_{AB}$  so that Alice can send Bob a message.
- They communicate with each other and the server as follows:

1. Alice sends the server S the names of Alice and Bob to request that a session key be generated.
  
2. The server sends to Alice:
  - a) The name of Bob
  - b) A session key for Alice and Bob to share
  - c) The name of Alice and the session key both encrypted using  $K_{BS}$

All 3 items above are encrypted using key  $K_{AS}$



3. Alice uses key  $K_{AS}$  to decrypt the items sent to her in step 2. Alice now knows the session key  $K_{AB}$ .
4. Alice sends Bob the value of 2c) which is the name of Alice and the session key  $K_{AB}$  encrypted with  $K_{BS}$
5. Bob decrypts the name of Alice and the session key using his key  $K_{BS}$ . Now Bob knows the session key  $K_{AB}$  which he uses to communicate with Alice.

# Needham-Schroeder

1.  $A \longrightarrow S: A, B$

2.  $S \longrightarrow A: e_{K_{AS}}(B, K_{AB}, e_{K_{BS}}(A, K_{AB}))$

Alice decrypts to get  $B, K_{AB}, e_{K_{BS}}(A, K_{AB})$

3.  $A \longrightarrow B: e_{K_{BS}}(A, K_{AB})$

$\longrightarrow$

Bob decrypts to get  $A, K_{AB}$

# Needham-Schroeder 2

1.  $A \longrightarrow S: A, B, N_A$
2.  $S \longrightarrow A: e_{K_{AS}}(B, N_A, K_{AB}, e_{K_{BS}}(A, K_{AB}))$
3.  $A \longrightarrow B: e_{K_{BS}}(A, K_{AB})$
4.  $B \longrightarrow A: e_{K_{AB}}(N_B)$
5.  $A \longrightarrow B: e_{K_{AB}}(N_B - 1)$

# Certificates

A certificate consists of a public key together with an identification of the key user. The certificate is issued by a trusted third party(TTP) called a

*certification agency (CA)*

The certification agency might be a government agency or financial institution.

The CA guarantees the link between the user and the public key by digitally signing a document which contains the user name, the public key, the name of the CA, the expiry date of the certificate and perhaps other information such as access rights.

# X.509 Standard

- Bob generates a document containing his relevant information and presents himself with this document to the CA.
- The CA confirm Bob's identity.
- The CA hash the document using SHA-1 and encrypt it using their own private key.
- This is the certificate.

- If Alice wants to communicate with Bob she looks up his public key document and certificate.
- She will use the public key of the CA to decrypt the certificate.
- She will hash the document using SHA-1
- If these two items are the same then she knows that she can safely communicate with Bob using the public key since the CA has verified his identity.